

Kenmerk: EWI2021/TW/MOR/MU/Mod7/Exam1

**Exam 1, Module 7, Codes 202001360 & 202001364**

**Discrete Structures & Efficient Algorithms**

Monday, March 25, 2022, 8:45 - 11:45

Answers to questions 1-8 need to be motivated, arguments and proofs must be complete. You are allowed to use a handwritten cheat sheet (A4, both sides) per topic (ADS, DM) during the exam.

For information: This exam consists of two parts, with the following (estimated) times per part:

Algorithms & Data Structures (ADS)	ca. 1h	(30 points)
Discrete Mathematics (DM)	ca. 2h	(60 points)

The total is  $30+60=90$  points. Your grade is  $1 + 0.1x$ ,  $x$  being the number of points, rounded to one digit. That means, you need 45 points to get a 5.5.

Please use a new sheet of paper for each part (ADS, DM), as the ADS and DM parts will be corrected separately!

Double students Discrete Mathematics & Algebra (202001364) only do the DM part. In that case, please write "Discrete Mathematics & Algebra" on top of your exam.

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# Algorithms & Data Structures

1. (10 points)

(a) Examine the following algorithm:

```
def func(n):
    res=0

    while n>0:
        m=n
        while m>0:
            res=res+m
            m=m-1
        n=n-1

    return res
```

Give the asymptotic time complexity of this algorithm expressed in the number of arithmetic operations.

(b) Suppose that the number of steps of an algorithm  $T(n)$  with an input of  $n$ , has the recurrence relation

$$T(n) = 8T(n/2) + n^2 + 4n + 1/n$$

What is the asymptotic complexity class of this algorithm?

2. (10 points)

(a) Give an algorithm that deletes the smallest element in a minheap, returning a heap with the remaining elements. The complexity of this algorithm must be  $O(\log n)$ .

(b) Given a non-empty binary tree sorted post-order. Give an algorithm that yields the node with the biggest element and the node with the smallest element.

3. (10 points)

Given  $n$  positive integers  $a_1, \dots, a_n$  and positive integer  $G$ . We want to determine whether there is a subsequence of  $a_1, \dots, a_n$  with sum  $G$ . Example: for 7, 3, 2, 5, 8 and  $G = 14$  the answer is "true" since  $7 + 2 + 5 = 14$ .

Define boolean  $R(i, g)$  as:

$R(i, g)$  iff  $a_1, \dots, a_i$  has subsequence with sum  $g$ .

Note that

- $R(i, 0) = \text{true}$  for all  $0 \leq i \leq n$

- $R(0, g) = \text{false}$  for all  $g > 0$
- $R(i, g) = \text{false}$  for all  $g < 0$

- (a) Motivate which of the following recurrence relations holds (for all  $1 \leq i \leq n$ ,  $1 \leq g \leq G$ ):
- $R(i, g) = (R(i-1, g) \leq a_i + R(i, g))$
  - $R(i, g) = R(i-1, g) \vee R(i-1, g - a_i)$
  - $R(i, g) = \min\{R(i+1, g), R(i+1, g + a_i)\}$
  - $R(i, g) = R(i-1, g) \vee R(i, g - a_i) \vee R(i-1, g)$
- (b) Give an algorithm to determine whether  $a_1, \dots, a_n$  has a subsequence with sum  $G$ . The complexity may be no worse than  $\Theta(Gn)$ .

**Reminder: Use a new sheet of paper for each part (ADS, DM), as the ADS and the DM parts will be corrected separately!**

## Discrete Mathematics

4. (4+5 points)

- (a) Consider  $a \in \mathbb{Z}$ , with  $a > 1$  and its (unique) prime factorization

$$a = p_1^{a_1} \cdot p_2^{a_2} \cdot \dots \cdot p_k^{a_k}.$$

Show that  $a$  is a perfect square<sup>1</sup> if and only if for all  $1 \leq i \leq k$ ,  $a_i$  is even.

- (b) Consider  $a, b \in \mathbb{Z}_{>0}$ , with  $\gcd(a, b) = 1$ . Show that if  $ab$  is a perfect square then both  $a$  and  $b$  must be perfect squares.

[Hint: You may want to use the statement of (a).]

5. (8 points) Consider a bipartite graph  $G = ((V_1, V_2), E)$  for which every vertex has the same positive degree. In other words, there exists an integer  $\Delta \geq 1$  such that  $\deg(v) = \Delta$  for all  $v \in (V_1 \cup V_2)$ . Prove that  $|V_1| = |V_2|$ .
6. (10 points) Consider a simple, capacitated network  $G = (V, E, c)$ , where  $V$  is the set of vertices,  $s, t \in V$ ,  $E$  is the set of directed edges, and  $c : E \rightarrow \mathbb{Z}_{\geq 0}$  gives the edge capacities. Let  $|V| = n$  and  $|E| = m$ . Suppose you are given a maximum  $(s, t)$ -flow  $f$ . Now suppose that the capacity of a specific edge  $e \in E$  is increased by one unit, and let  $G'$  be the resulting capacitated network. Suggest how to compute the maximum  $(s, t)$ -flow in  $f'$  in  $G'$ , in  $O(n+m)$  time. Briefly explain (i) why your suggested algorithm

<sup>1</sup>An  $x \in \mathbb{Z}$  is a perfect square if and only if there exists  $y \in \mathbb{Z}$  such that  $x = y^2$ .

is correct (ii) why it achieves the desired running time.

[Hint: You may want to use the  $(s, t)$ -flow  $f$  as well as the residual graph for  $G'$  with respect to  $f$ .]

7. (4+3 points)

(a) Compute the solution to the recurrence relation

$$a_n - 7a_{n-1} + 10a_{n-2} = 6 + 8n \quad (n \geq 2) \quad \text{with} \quad a_0 = 1 \text{ and } a_1 = 2.$$

(b) Let  $a_n$  with  $n \geq 0$  be the number of bit strings (i.e., strings in  $\{0, 1\}$ ) of length  $n$  that contain three consecutive zeros. Find a recurrence relation for computing  $a_n$ . (You do not need to solve this recurrence relation.)

8. (8 points) Assume that Alice has published modulus  $n = 55$ , and exponent  $e = 7$ . Bob sends ciphertext  $C = 5$  to Alice. You are eavesdropper Eve and you are interested in Bob's secret message  $M$ . Compute Bob's secret message  $M$  from ciphertext  $C$ . Write down all of the computational steps that you need to perform in order to obtain Bob's secret message  $M$ .

9. (3 points each) For each of the following six claims, decide if true or false or if you would rather not give an answer. A correct answer gives **3**, an incorrect answer **-3** and not giving an answer **0 points** (minimum total number of points for Question 9 is 0 points). **Instead of guessing, it may be better not giving an answer.**

(a) Consider an undirected, simple graph  $G = (V, E)$  with edge weights  $w_e \geq 0$ ,  $e \in E$ . Consider a cycle  $C$  in  $G$ , and let  $e$  be an edge of the cycle  $C$  such that  $w_e \leq w_{e'}$  for any edge  $e'$  along  $C$ . Then  $e$  must belong to a minimum spanning tree of  $G$ .

True.....

False.....

I prefer to not give an answer.....

(b) Consider a capacitated network  $G = (V, A, c)$ , where  $V$  is the set of vertices,  $A$  is the set of directed arcs, and  $c : A \rightarrow \mathbb{Z}_{\geq 0}$ , are the arc capacities. Let  $f$  be some  $(s, t)$ -flow in  $G$  respecting the flow conservation and capacity constraints, and let  $val(f)$  be its value. Then there must exist an  $(s, t)$ -cut  $[S, T]$  with capacity  $cap([S, T]) = val(f)$ .

True.....

False.....

I prefer to not give an answer.....

- (c) Consider an undirected, simple graph  $G = (V, E)$  with edge weights  $w : E \rightarrow \mathbb{R}$ , and let the minimum edge weight be  $-10$  (i.e.,  $\min_e \{w(e)\} = -10$ ). The following algorithm can be used to compute the shortest path between  $u, v \in V$  in  $G$ : (i) Construct graph  $G' = (V, E)$  with edge weights  $w' : E \rightarrow \mathbb{R}_{\geq 0}$  defined as  $w'(e) = w(e) + 10$ , for all  $e \in E$ . (ii) Run Dijkstra's algorithm to compute the shortest path  $P(u, v)$  between  $u, v$  in  $G'$ . (Note that by construction  $G'$  only has non-negative weight edges). (iii) Output  $P(u, v)$  as the shortest path between  $u, v$  in  $G$ .

True .....   
False .....   
I prefer to not give an answer .....

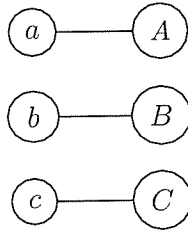
- (d) Consider an undirected, connected, and simple graph  $G = (V, E)$  with non-negative edge weights. Also consider a minimum spanning tree  $T$  of  $G$ . Then for every pair of vertices  $u, w \in V$ , the shortest  $(u, w)$ -path in  $G$  is contained in  $T$ .

True .....   
False .....   
I prefer to not give an answer .....

- (e) Consider a connected graph  $G = (V, E)$  with 13 vertices (i.e.,  $|V| = 13$ ) such that  $G$  has an Euler path but no Euler tour. Then  $G$  must have exactly two vertices of odd degree and eleven vertices of even degree.

True .....   
False .....   
I prefer to not give an answer .....

(f) Consider the depicted matching  $M$  and corresponding preference lists. Then  $M$  is a stable matching.



$B >_a A >_a C$   
 $C >_b B >_b A$   
 $A >_c C >_c B$   
 $b >_A a >_A c$   
 $c >_B b >_B a$   
 $a >_C c >_C b$

[Reminder:  $x >_z y$  indicates that  $z$  prefers to be matched with  $x$  over  $y$ .]

- True .....   
 False .....   
 I prefer to not give an answer .....