

# Test Pearl 100 — Cryptography

Pearls of Computer Science (201300070)  
6 October 2017, 13:45–14:45  
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- You may use 1 A4 sheet with your own notes for this test, as well as a *simple* calculator
- Scientific or graphical calculators, laptops, mobile phones, books etc. are not allowed.  
*Put those in your bag now (with the sound switched off)!*
- Always motivate/explain your answers, unless it is explicitly stated not to!
- Total number of points: 31

## 12 points Question 1

- Encrypt the message CRYPTO with key EASILY using the Vigenère cipher.
- What is the main practical drawback of the One-Time-Pad?
- Consider the following keyed-function

$$F : \{0, 1\}^3 \times \{0, 1\}^3 \rightarrow \{0, 1\}^3 \quad \text{whereas} \quad F(K, x) = K \oplus x.$$

Encrypt the 6-bit message  $m = 110011$  using the Feistel cipher with 3 rounds, above keyed-function  $F$ , and the round-keys  $K_0 = 111$ ,  $K_1 = 010$ , and  $K_2 = 011$ .

- Consider the following plaintext message (a 4-bit string)

1001

Encrypt this message in the CBC-mode by using the following 2-bit block cipher

$$E_k(b_1b_0) = b_1b_0 \oplus k$$

with the bit-string  $k = 01$  as secret key (note that  $b_1b_0$  denotes an arbitrary 2-bit plaintext message). As initialization vector for the CBC-mode, use the bit-string  $IV = 11$ .

## 7 points Question 2

(Hint. Do not try to solve any of the following by using the efficient exponentiation modulo  $N$ .)

- Write down all the elements in  $\mathbb{Z}_{16}^*$ .
- Compute  $(25^{2212122} - 2) \bmod 13$  using modular arithmetic (recall that the result needs to be  $\geq 0$ ).
- Let  $p = 43$  and  $q = 41$  be primes, and  $N = pq = 1763$  an RSA-modulus. Compute  $1326^{1680} \bmod 1763$ .

## 8 points Question 3

Let  $p = 37$ ,  $q = 31$ , and  $N = pq = 1147$ . Assume that we use  $(N, e) = (1147, 463)$  as the public key in the RSA signature scheme.

- Compute Euler's totient function  $\phi(N)$ .
- Use the extended Euclidean algorithm (it is mandatory to use the table method here!) to compute the secret key  $d \geq 0$  that corresponds to the public key  $(N, e) = (1147, 463)$ .
- Sign the message  $m = 2$  using the RSA signature scheme with the in (b) computed secret key  $d \geq 0$  (if you couldn't solve (b), then use the key  $d = 11$ , which is different from the correct result in (b)!).

## 4 points Question 4

Let  $(N, e_1)$  and  $(N, e_2)$  be two RSA public keys with the same RSA-modulus  $N$  but different public exponents  $e_1 \neq e_2$  such that  $\gcd(e_1, e_2) = 1$ . Suppose that you are given the following two ciphertexts of the same message  $m$  (the message  $m$  is not known to you):

$$c_1 = m^{e_1} \bmod N \quad \text{and} \quad c_2 = m^{e_2} \bmod N.$$

Provide a way to get to the underlying plaintext  $m$  by finding integers  $x$  and  $y$  such that

$$c_1^x \cdot c_2^y \bmod N = m.$$