

Test Pearl 100 — Cryptography

Pearls of Computer Science (201300070)

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Module coordinator: Doina Bucur

Instructor: Andreas Peter

- You may use 1 A4 sheet with your own notes for this test, as well as a *simple* calculator
- Scientific or graphical calculators, laptops, mobile phones, books etc. are not allowed.
Put those in your bag now (with the sound switched off)!
- Always motivate/explain your answers, unless it is explicitly stated not to!
- Total number of points: 38

14 points **Question 1** Write down your answers (A, B, C, or D) on your answer sheet and **NOT** on the exam sheet at hand. There is only **one correct answer per subquestion**. Your answers require no motivation.

IMPORTANT: Don't just guess; for each wrong answer, you get 1 point deducted!

- (a) Let JVVIMQ be a ciphertext produced by the Vigenère cipher using the key PIN. Which of the following is the underlying plaintext?
- ATTACK
 - DEFEND
 - UNITED
 - CRYPTO
- (b) According to the second Kerckhoffs' principle, which of the following statements is correct?
- The secret key used in a secret-key encryption scheme must be able to fall into the hands of the enemy without inconvenience.
 - The public key used in a public-key encryption scheme must be able to fall into the hands of the enemy without inconvenience.
 - The source code of the implementation of an encryption scheme must **not** be able to fall into the hands of the enemy.
 - The output of an encryption scheme (i.e., a ciphertext) must **not** be able to fall into the hands of the enemy.
- (c) Suppose that Alice encrypts the plaintext 000 using the One-Time-Pad. Assuming that you don't know which key Alice used in the encryption, what is the probability that 000 is the resulting ciphertext?
- 0%
 - 6.25%
 - 12.5%
 - 25%
- (d) What is the concept of "hybrid encryption"?
- First use a secret-key encryption scheme to exchange a key, then use this key in a public-key encryption scheme.
 - First encrypt a given message with a secret-key encryption scheme, then produce a public-key signature on the resulting ciphertext.
 - First use a public-key encryption scheme to exchange a key, then use this key in a secret-key encryption scheme.
 - First produce a public-key signature on a given message, then encrypt the resulting signature with a secret-key encryption scheme.
- (e) Let $p = 31$ and $q = 43$ be primes, and $N = pq = 1333$. What is the result of the computation: $(1234^{1260} + 2567) \bmod 1333$?
- 1232
 - 1233
 - 1234
 - 1235

- (f) Which elements are contained in \mathbb{Z}_{12}^* ?
 - A. 1, 5, 7, 11
 - B. 0, 1, 5, 7
 - C. 3, 5, 7, 11
 - D. 1, 3, 5, 7
- (g) Which of the following numbers is a valid (= generated as described in the lecture), but too small to be secure, RSA modulus?
 - A. $N = 9$
 - B. $N = 29$
 - C. $N = 1024$
 - D. $N = 1271$

7 points **Question 2** Consider the following keyed-function

$$F : \{0, 1\}^2 \times \{0, 1\}^2 \rightarrow \{0, 1\}^2 \quad \text{whereas} \quad F(K, x) = K \oplus x.$$

Decrypt the 4-bit ciphertext $c = 1101$ using the Feistel cipher with 3 rounds, above keyed-function F , and the round-keys $K_0 = 11$, $K_1 = 10$, and $K_2 = 01$.

9 points **Question 3** Consider the following plaintext message (a 5-bit string)

11010

Encrypt this message in the CBC-mode by using the following 2-bit block cipher

$$E_k(b_1b_0) = b_1b_0 \oplus k$$

with the bit-string $k = 10$ as secret key (note that b_1b_0 denotes an arbitrary 2-bit plaintext message). As initialization vector for the CBC-mode, use the bit-string $IV = 11$.

8 points **Question 4** Let $p = 37$, $q = 41$, and $N = pq = 1517$. Assume that we use $(N, e) = (1517, 823)$ as the public key in the RSA signature scheme.

- (a) Compute Euler's totient function $\phi(N)$.
- (b) Compute the RSA secret key $d \geq 0$ that corresponds to the public key $(N, e) = (1517, 823)$. The following table with some pre-computed calculations might be helpful (as it avoids the need to compute the extended Euclidean algorithm):

A. $\gcd(p, q) = 1$	E. $1 = 1440 \cdot 197 + 1517 \cdot (-187)$
B. $\gcd(\phi(N), N) = 1$	F. $1 = 37 \cdot 10 + 41 \cdot (-9)$
C. $\gcd(\phi(N), e) = 1$	G. $1 = 1517 \cdot (-319) + 823 \cdot 588$
D. $\gcd(N, e) = 1$	H. $1 = 1440 \cdot (-4) + 823 \cdot 7$

- (c) Sign the message $m = 2$ using the RSA signature scheme with the in (b) computed secret key $d \geq 0$ (if you couldn't solve (b), then use the key $d = 11$, which is different from the correct result in (b)!).